# Radio Coloring Phenomena and Its Applications 

B.R. Srinivas<br>Associate Professor of Mathematics, St. Marys Group of Institutions<br>Guntur, A.P, INDIA<br>brsmastan@gmail.com

A. Sri Krishna Chaitanya<br>Associate Professor of Mathematics, Chebrolu Engineering College, Guntur. A.P, INDIA askc_7@yahoo.com

Abstract: This paper studies the concepts of Radio Colorings. The main results are<br>1) Relation between radio chromatic number for a connected graph of order $n$ having diameter d, and for integers $k, l$.<br>2) Bounds for radio number of a connected graph of order $n$ and diameter $d$.<br>Mathematics Subject Classification 2000: 05CXX,05C15,05C20,37E25.<br>Key words: connected graph, positive integers, and chromatic number.

## 1. Introduction

The concept of $\mathrm{L}(\mathrm{h}, \mathrm{k})$-colorings has been generalized in a natural way. For nonnegative integers $\mathrm{dl}, \mathrm{d} 2, \ldots, \mathrm{dk}$, where $\mathrm{k} \geq 2$, and $\mathrm{L}(\mathrm{dl}, \mathrm{d} 2, \ldots, \mathrm{dk})$-coloring c of a graph G is an assignment c of colors (nonnegative integers in this case) to the vertices of G such that $|c(u)-c(w)| \geq \mathrm{di}$ whenever $\mathrm{d}(\mathrm{u}, \mathrm{w})=\mathrm{i}$ for $1 \leq \mathrm{i} \leq \mathrm{k}$. The $\mathrm{L}(\mathrm{d} 1, \mathrm{~d} 2, \ldots, \mathrm{dk})$-colorings in which $\mathrm{di}=\mathrm{k}+1-\mathrm{i}$ for each i ( $1 \leq \mathrm{i} \leq \mathrm{k}$ ) have proved to be of special interest.
The term "radio coloring" emanates from its connection with the Channel Assignment Problem. In the United States, one of the responsibilities of the Federal Communications Commission (FCC) concerns the regulation of FM radio stations. Each station is characterized by its transmission frequency, effective radiated power, and antenna height. Each FM station is assigned a station class, which depends on a number of factors, including its effective radiated power and antenna height. The FCC requires that FM radio stations located within a certain proximity to one another must be assigned distinct channels and that the nearer two stations are connected to each other, the greater the difference in their assigned channels must be minimum distance between stations (see [6]).
We have also mentioned that the use of graph theory to study the Channel Assignment Problem and related problems dates back at least to 1970(Metzger [5]) .In 1980, William Hale [2] modeled the Channel Assignment Problem as both a frequency-distance constrained and frequency constrained optimization problem and discussed applications to important real world problems.

## 2. Preliminaries

### 2.1.Definition:

A radio coloring of $G$ is an assignment of colors to the vertices of $G$ such that two colors $i$ and $j$ can be assigned to two distinct vertices $u$ and $v$ only if $d(u, v)+|i-j| \geq 1+k$ for some fixed positive integer $k$.

### 2.2.Definition:

For a connected graph $G$ of diameter $d$ and an integer $k$ with $1 \leq k \leq d$, a $k$-radio coloring $c$ of $G$ (sometimes called a radio $k$-coloring) is an assignment of colors (positive integers) to the vertices of $G$ such that

$$
d(u, v)+|c(u)-c(v)| \geq 1+k
$$

for every two distinct vertices $u$ and $v$ of $G$. Thus a 1 -radio coloring of $G$ is simply a proper coloring of $G$, while a 2 -radio coloring is an $L(2,1)$-coloring [1]. Note that a $k$-radio coloring $c$ does not imply that $c$ is a $k$-coloring of the vertices of $G$ (a vertex coloring using $k$ colors).

### 2.3. Definition

The value $r c_{k}(c)$ of a $k$-radio coloring $c$ of $G$ is defined as the maximum color assigned to a vertex of $G$ by $c$ (where, again, we may assume that some vertex of $G$ is assigned the color $1)$. The coloring $\bar{c}$ of $G$ defined by

$$
\bar{c}(v)=r c_{k}(c)+l-c(v)
$$

for every vertex $v$ of $G$ is also a $k$-radio coloring of $G$, referred to as the complementary coloring of $c$. Because it is assumed that some vertex of $G$ has been colored 1 by $c$, it follows that $r c_{k}(\bar{c})=r c_{k}(c)$.

### 2.4.Definition

For a connected graph $G$ with diameter $d$ and an integer $k$ with $1 \leq k \leq d$, the $k$-radio chromatic number (or simply the $k$-radio number) $r c_{k}(G)$ is defined as

$$
r c_{k}(G)=\min \left\{r c_{k}(c)\right\}
$$

Where the minimum is taken over all $k$-radio colorings $c$ of $G$. Since a 1 -radio coloring of $G$ is a proper coloring, it follows that $r c_{l}(G)=\chi(G)$. On the other hand, a 2-radio coloring of $G$ is an $L(2,1)$-coloring of $G$, all of whose colors are positive integers. Thus

$$
r c_{2}(G)=1+\lambda(G)
$$

## 3. Radio Chromatic Number:

### 3.1.Proposition:

For a connected graph $G$ of order $n$ having diameter $d$ and for integers $k$ and $\ell$ with $l \leq k<\ell \leq$ $d$.

$$
r c_{l}(G) \leq r c_{k}(G)+(n-1)(\ell-k) .
$$

Proof. Let c be a $k$-radio coloring of $G$ such that $r c_{k}(c)=r c_{k}(G)$. Let $V(G)=\left\{v_{l}, v_{2}, \ldots, v_{n}\right\}$ such that $c\left(v_{i}\right) \leq c\left(v_{i+1}\right)$ for $1 \leq i \leq n-1$. We define a coloring $c^{\prime}$ of $G$ by

$$
c^{\prime}\left(v_{i}\right)=c\left(v_{i}\right)+(i-1)(l-k)
$$

For integers $i$ and $j$ with $l \leq i<j \leq n$, we therefore have

$$
\left|c^{\prime}\left(v_{i}\right)-c^{\prime}\left(v_{j}\right)\right|=\left|c\left(v_{i}\right)-c\left(v_{j}\right)\right|+(j-i)(l-k)
$$

Since $c$ is a $k$-radio coloring of $G$, it follows that

$$
\left|c\left(v_{i}\right)-c\left(v_{j}\right)\right| \geq 1+k-d\left(v_{i}, v_{j}\right)
$$

Consequently,

$$
\begin{aligned}
\left|c^{\prime}\left(v_{i}\right)-c^{\prime}\left(v_{j}\right)\right| & \geq 1+k+(j-i)(l-k)-d\left(v_{i}, v_{j}\right) \\
& \geq 1+l-d\left(v_{i}, v_{j}\right)
\end{aligned}
$$

Thus $c^{\prime}$ is an $l$-radio coloring of $G$ with

$$
r c_{l}\left(c^{\prime}\right)=r c_{k}(G)+(n-1)(\ell-k)
$$

and so $r c_{l}(G) \leq r c_{k}(G)+(n-1)(\ell-k)$.

### 3.2.Definition:

Even though $k$-radio colorings of a connected graph with diameter $d$ are defined for every integer $k$ with $1 \leq k \leq d$, it is the two smallest and two largest values of $k$ that have received the most attention. For a connected graph $G$ with diameter $d$, a $d$-radio coloring $c$ of a connected graph $G$ with diameter $d$ requires that

$$
d(u, v)+|c(u)-c(v)| \geq 1+d
$$

for every two distinct vertices $u$ and $v$ of $G$. A $d$-radio coloring is called a radio labeling and the $d$-radio chromatic number (or $d$-radio number) is sometimes called simply the
radio number $r n(G)$ of $G$.

### 3.3.Proposition

If $G$ is a connected graph of order $n$ and diameter $d$, then

$$
n \leq r n(G) \leq 1+(n-1) d
$$

We have noted that if $d=1$ and so $G=K_{n}$, then $r n\left(K_{n}\right)=n$. The graph $C_{5}$ and the Petersen graph $P$ both have diameter 2 and their radio numbers also attain the lower bound, namely $r n\left(C_{5}\right)=5$ and $\quad r n(P)=10$. Furthermore, for each integer $k \geq 2$, the graph $K_{k} \times K_{2}$ has order $n=2 k$, diameter 2, and
$r n\left(K_{k} \times K_{2}\right)=n$. The graph $C_{3} \times C_{5}$ has order $\mathrm{n}=15$, diameter 3, and radio number 15.(Figure 1)


Figure1. A radio labeling of $C_{3} \times C_{5}$
For a connected graph $G$ of order $n$ and diameter $d$, the upper bound for $r n(G)$ given in 2.3 can often be improved.

### 3.4.Proposition:

If $G$ is connected graph of order $n$ and diameter $d$ containing an induced sub graph $H$ of order $p$ and diameter $d$ such that $d_{H}(u, v)=d_{G}(u, v)$ for every two vertices $u$ and $v$ of $H$, then

$$
r n(H) \leq r n(G) \leq r n(H)+(n-p) d .
$$

A special case of Proposition 2.4 is when $H$ is a path.

### 3.5.Corollary

If $G$ is a connected graph of order $n$ and diameter d ,
Then $r n\left(P_{d+1}\right) \leq r n(G) \leq r n\left(P_{d+1}\right)+(n-d-1) d$
Corollary 2.5 illustrates the value of knowing the radio numbers of paths. The following result was obtained by Daphne Liu and Xuding Zhu [4].

### 3.6.Theorem:

For every integer $n \geq 3$,

$$
r n\left(P_{n}\right)=\left\{\begin{array}{l}
2 r^{2}+3 \quad \text { if } n=2 r+1 \\
2 r^{2}-2 r+2 \text { if } n=2 r
\end{array}\right.
$$

Combining Corollary 2.5 and Theorem2.6, we have the following.

### 3.7.Corollary:

Let $G$ be a connected graph of order $n$ and diameter $d$.
a) If $d=2$, then $4 \leq r n(G) \leq 2 n-2$.
b) If $d=3$, then $6 \leq r n(G) \leq 3 n-6$.
c) If $d=4$, then $11 \leq r n(G) \leq 4 n-9$.

## Proof:

While the paths $P_{d+1}$ show the sharpness of the lower bounds in Corollary 2.7 the sharpness of the upper bounds are less obvious. It is not difficult to show that for every integer $n \geq 3$, there exists a connected graph $G$ of diameter 2 with $r n(G)=2 n-2$. The graph $H$ of Figure 2(a) has order $n=6$,
$\operatorname{diam}(H)=3$ and $\mathrm{rn}(H)=12=3 n-6$. The
graph $F$ of Figure 2(b) has order $n=6$, $\operatorname{diam}(F)=4$, and $\quad \operatorname{rn}(F)=14=4 n-10$. The number $4 n-10$ does not attain the upper bound for the radio number of a graph of diameter 4 given in Corollary 2.7(c). In fact, it may be that the appropriate upper bound for this case is $4 n-10$ rather than $4 n-9$.


Figure2. Radio numbers of graphs having diameters 3 \& 4
For a connected graph $G$ of diameter $d$, a $(d-1)$-radio coloring c requires that

$$
\mathrm{d}(\mathrm{u}, \mathrm{v})+|c(u)-c(v)| \geq \mathrm{d}
$$

for every two distinct vertices $u$ and $v$ of $G$. A $(d-1)$-radio coloring $c$ is also referred to as a radio antipodal coloring (or simply an antipodal coloring) of $G$ since $c(u)=c(v)$ only if $u$ and $v$ are antipodal vertices of $G$. The radio antipodal number or, more simply, the antipodal number $a n(c)$ of $c$ is the largest color assigned to a vertex of $G$ by $c$. The antipodal chromatic number or the antipodal number $a n(G)$ of $G$ is

$$
a n(G)=\min \{a n(c)\}
$$

where the minimum is taken over all radio antipodal colorings $c$ of $G$. If $c$ is a radio antipodal coloring of a graph $G$ such that $a n(c)=\ell$, then the complementary coloring $\bar{c}$ of $G$ defined by

$$
\bar{c}(v)=\ell+1-c(v)
$$

for every vertex $v$ of $G$ is also a radio antipodal coloring of $G$.
A radio antipodal coloring of the graph $H$ in Figure 3 is given with antipodal number 5. Thus $a n(H) \leq 5$. Let c be a radio antipodal coloring of $H$ with $a n(c)=a n(H) \leq 5$. Since $\operatorname{diam}(H)=3$, the colors of every two adjacent vertices of $H$ must differ by at least 2 and the colors of two vertices at distance 2 must differ. We may assume that $c(v) \in\{1,2\}$, for otherwise, $\bar{c}(v) \in\{1,2\}$ for the complementary radio antipodal coloring $\bar{c}$ of $H$.
Suppose that $c(v)=a \leq 2$. Then at least one of the vertices $u, w$, and $y$ must have color at least $\quad a+2$, one must have color at least $a+3$, and the other must have color at least $a+4$. Since $a+4 \geq 5$, it follows that $a n(c) \geq 5$ and so $a n(H) \geq 5$. Hence $a n(H)=5$



Figure 3. A graph with antipodal number 5


Figure 4. Radio antipodal colorings of $P_{n}(3 \leq n \leq 6)$
Figure 4 gives radio antipodal colorings of the paths $P_{n}$ with $3 \leq n \leq 6$ that give $a n\left(P_{n}\right)$ for these graphs. The antipodal numbers of all paths were determined by Khennoufa and Togni [3].

### 3.8.Remarks:

1) If we consider less diameter, Radio coloring concept vanishes.
2) If we consider maximum diameter, required radio chromatic number will occur.
3) The conditions of radio coloring is not applicable for chromatic polynomials.

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## Authors' Biography



Mr. B.R. Srinivas, completed his M.Sc., AO Mathematics from Acharya Nagarjuna university in the year 1995, M.Phil Mathematics from Madurai Kamaraj University in the year 2005 and M.Tech computer science engineering from Vinayaka Mission University in the year 2007. He attended for five international conferences, six faculty development programs and twenty national workshops. At present he is working as Associate Professor, St. Marys Group of Institutions Guntur, A.P, INDIA. His area of interest is graph theory and theoretical computer science.


Mr. A. Sri Krishna Chaitanya, completed his M.Sc., Mathematics from Acharya Nagarjuna university in the year 2005 and M.Phil Mathematics from Alagappa University in the year 2007. At present he is working as Associate Professor, Chebrolu Engineering College, Chebrolu, Guntur Dist, A.P, INDIA. He is interested to work in the areas of Graph Theory, Boolean algebra, Lattice Theory and the Related Fields of Algebra.

