# Representation of Pre A* - Algebra by a Partially Order 

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#### Abstract

This manuscript is a study on Pre-A*-algebra A in view of it is like a partially ordered set. Using a binary operation in Pre-A*-algebra, an observation is made on Pre $A^{*}$-Algebra as a partially ordered set with respect to binary operation $\wedge$ and obtained corresponding results. It is also make available an equivalent condition for a Pre $A^{*}$-algebra become a Boolean algebra.


Keywords: A*-algebra, Pre-A*-algebra, Boolean algebra, Partially ordered set, Ada, Homomorphism.
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## 1. Introduction

In a draft manuscript entitled "The Equational theory of Disjoint Alternatives", E. G. Manes (1989) introduced the concept of Ada (Algebra of disjoint alternatives) $\left(A, \wedge, \nu,(-)^{\prime},(-)_{\pi}, 0,1,2\right)$ which is however differs from the definition of the Ada of E. G. Manes (1993) later paper entitled "Adas and the equational theory of if-then-else". While the Ada of the earlier draft seems to be based on extending the If-Then-Else concept more on the basis of Boolean algebras and the later concept is based on C-algebras A ( $\wedge, \vee ‘)$ introduced by Fernando Guzman and Craig C. Squir (1990). P. Koteswara Rao (1994) first introduced the concept of A*-algebra $\left(A, \wedge, \vee, *,(-)^{\sim}(-)_{\pi}, 0,1,2\right)$ not only studied the equivalence with Ada, C-algebra, Ada's connection with 3-Ring, Stone type representation but also introduced the concept of A*-clone, the If-ThenElse structure over A*-algebra and Ideal of A*-algebra.
J.Venkateswara Rao (2000) introduced the concept Pre A*-algebra $\left(A, \wedge, \vee,(-)^{\sim}\right)$ analogous to Calgebra as a reduct of A*- algebra. Venkateswara Rao.J, Praroopa.Y (2006) made a structural study on Boolean algebras and Pre A*-Algebras.
Boolean algebra depends on two element logic. C-algebra, Ada, A*- algebra and our Pre A*algebra are regular extensions of Boolean logic to 3 truth values, where the third truth value stands for an undefined truth value. The Pre A*- algebra structure is denoted by $\left(A, \wedge, \vee,(-)^{\sim}\right)$ where A is non-empty set, $\wedge, \vee$ are binary operations and $(-)^{\sim}$ is a unary operation.

In this paper we define a relation $\leq$ on Pre $\mathrm{A}^{*}$-algebra with respect to the binary operation $\wedge$, we discuss the properties of a Pre A*-algebra like a poset. We find the necessary conditions for a poset to become a lattice. We also present a equivalent condition for a Pre $\mathrm{A}^{*}$-algebra become a Boolean algebra. For any $\mathrm{a} \in \mathrm{A}$ define $A_{a}=\{\mathrm{x} \in \mathrm{A} / \mathrm{a} \wedge \mathrm{x}=\mathrm{x}\}$ and $x^{a}=\mathrm{a} \wedge \mathrm{x}^{\sim}$ then $\left(A_{a}, \wedge, \vee,{ }^{\mathrm{a}}\right)$ OARC
is a Pre $\mathrm{A}^{*}$-algebra. We also define a mapping $\alpha_{a, b}$ from $A_{b}$ to $A_{a}$ by $\alpha_{a, b}(\mathrm{x})=\mathrm{a} \wedge \mathrm{x}$ for all $\mathrm{x} \in A_{b}$ is a homomorphism of $\operatorname{Pre} \mathrm{A}^{*}$-algebras.

## Preliminaries

1.1. Definition: The relation $R$ on a set $A$ is called a partial order on $A$ when $R(\leq)$ is reflexive, anti-symmetric, and transitive. Under these conditions, the set A is called a partially ordered set or a poset. Frequently we write $(\mathrm{A}, \mathrm{R})$ or $(\mathrm{A}, \leq)$ to denote that A is partially ordered by the relation $\mathrm{R}(\leq)$. Since the relation $\leq$ on the set of real numbers is the prototype of a partial order it is common to write $\leq$ to represent an arbitrary partial order can be described as follows:

1. For all $\mathrm{a} \in \mathrm{A}, \mathrm{a} \leq \mathrm{a} \quad$ (reflexive)
2. For all $\mathrm{a}, \mathrm{b} \in \mathrm{A}, \mathrm{a} \leq \mathrm{b}, \mathrm{b} \leq \mathrm{a}$, then $\mathrm{a}=\mathrm{b} \quad$ (anti symmetry)
3. For all $\mathrm{a}, \mathrm{b}, \mathrm{c} \in \mathrm{A}, \mathrm{a} \leq \mathrm{b}$ and $\mathrm{b} \leq \mathrm{c}$, then $\mathrm{a} \leq \mathrm{c} \quad$ (transitivity)

Two elements $a$ and $b$ in $A$ are said to be comparable under $\leq$ if either $\mathrm{a} \leq \mathrm{b}$ or $\mathrm{b} \leq \mathrm{a}$; otherwise they are incomparable. If every pair of elements of A are comparable, then we say that the partially ordered set is totally ordered.
1.2. Definition: An algebra $\left(A, \wedge, \vee,(-)^{\sim}\right)$ where A is a non-empty set with $1, \wedge, \vee$ are binary operations and $(-)^{\sim}$ is a unary operation satisfying
(a) $x^{\sim}=x \quad \forall x \in A$
(b) $x \wedge x=x, \quad \forall x \in A$
(c) $x \wedge y=y \wedge x, \quad \forall x, y \in A$
(d) $(x \wedge y)^{\sim}=x^{\sim} \vee y^{\sim} \quad \forall x, y \in A$
(e) $x \wedge(y \wedge z)=(x \wedge y) \wedge z, \quad \forall x, y, z \in A$
(f) $x \wedge(y \vee z)=(x \wedge y) \vee(x \wedge z), \quad \forall x, y, z \in A$
(g) $x \wedge y=x \wedge\left(x^{\sim} \vee y\right), \quad \forall x, y \in A$ is called a Pre $A^{*}$-algebra.
1.1. Example: $\mathbf{3}=\{0,1,2\}$ with operations $\wedge, \vee,(-)^{\sim}$ defined below is a Pre $A^{*}$-algebra.

| $\wedge$ | 0 | 1 | 2 |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 0 | 2 |
| 1 | 0 | 1 | 2 |
| 2 | 2 | 2 | 2 |


| $\vee$ | 0 | 1 | 2 |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 2 |
| 1 | 1 | 1 | 2 |
| 2 | 2 | 2 | 2 |


| $x$ | $x^{\sim}$ |
| :--- | :--- |
| 0 | 1 |
| 1 | 0 |
| 2 | 2 |

1.1. Note: The elements $0,1,2$ in the above example satisfy the following laws:
(a) $2^{\sim}=2$
(b) $1 \wedge x=x$ for all $x \in \mathbf{3}$
(c) $0 \vee \mathrm{x}=\mathrm{x}$ for all $\mathrm{x} \in \mathbf{3}$
(d) $2 \wedge x=2 \vee x=2$ for all $x \in 3$.
1.2. Example: $\mathbf{2}=\{0,1\}$ with operations $\wedge, \vee(-)^{\sim}$ defined below is a Pre $A^{*}$-algebra.

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| $\wedge$ | 0 | 1 |  | $\vee$ | 0 | 1 |  | $x$ | $x^{\sim}$ |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |
|  | 0 | 0 |  | 0 | 1 |  | 0 | 1 |  |
| 1 | 0 | 1 |  | 1 | 1 | 1 |  | 1 | 0 |

1.2. Note :(i) $(2, \vee, \wedge,(-))$ is a Boolean algebra. So every Boolean algebra is a Pre $A^{*}$ algebra.
(ii) The identities 1.2 (a) and $1.2(\mathrm{~d})$ imply that the varieties of Pre $\mathrm{A}^{*}$-algebras satisfies all the dual statements of $1.2(\mathrm{~b})$ to $1.2(\mathrm{~g})$.
1.3. Definition: Let $A$ be a Pre $A^{*}$-algebra. An element $x \in A$ is called a central element of $A$ if $x \vee x^{\sim}=1$ and the set $\left\{\mathrm{x} \in \mathrm{A} / x \vee x^{\sim}=1\right\}$ of all central elements of A is called the centre of A and it is denoted by $\mathrm{B}(\mathrm{A})$.
1.1. Theorem: [Satyanarayana.A, (2012)] Let $A$ be a Pre $A^{*}$-algebra with 1, then $B(A)$ is a Boolean algebra with the induced operations $\wedge, \vee,(-)^{\sim}$
1.1. Lemma: [Satyanarayana.A, (2012)] Every Pre $A^{*}$-algebra with 1 satisfies the following laws
(a) $x \vee 1=x \vee x^{\sim}$
(b) $x \wedge 0=x \wedge x^{\sim}$
1.2. Lemma: [Satyanarayana.A, (2012)] Every Pre $A^{*}$-algebra with 1 satisfies the following laws.
(a) $x \wedge\left(x^{\sim} \vee x\right)=x \vee\left(x^{\sim} \wedge x\right)=x$
(b) $\left(x \vee x^{\sim}\right) \wedge y=(x \wedge y) \vee\left(x^{\sim} \wedge y\right)$
(c) $(x \vee y) \wedge z=(x \wedge z) \vee\left(x^{\sim} \wedge \mathrm{y} \wedge \mathrm{z}\right)$

1. 4. Definition: Let $\left(A_{1}, \vee, \wedge,(-)^{\sim}\right)$ and $\left(A_{2}, \vee, \wedge,(-)^{\sim}\right)$ be a two Pre $\mathrm{A}^{*}$ - algebras. A mapping $f: A_{1} \rightarrow A_{2}$ is called a Pre $\mathrm{A}^{*}$-homomorphism if
(i) $f(a \wedge b)=f(a) \wedge f(b)$
(ii) $f(a \vee b)=f(a) \vee f(b)$
(iii) $f\left(a^{\sim}\right)=(f(a))^{\sim}$

The homomorphism $f: A_{1} \rightarrow A_{2}$ is onto, then f is called epimorphism.
The homomorphism $f: A_{1} \rightarrow A_{2}$ is one-one then f is called monomorphism
The homomorphism $f: A_{1} \rightarrow A_{2}$ is one-one and onto then $f$ is called an isomorphism, and $A_{1}, A_{2}$ are isomorphic, denoted in symbol $A_{1} \cong A_{2}$.

## 2. Pre A*- Algebra as a Poset with Respect to Binary Operation $\wedge$

2. 1 Definition: Let $A$ be a Pre $A^{*}$-algebra. Define a relation $\leq o n A$ by $x \leq y$ if and only if $y \wedge x$ $=x \wedge y=x$.
3. 1 Lemma: If A is a Pre $\mathrm{A}^{*}$-algebra, then $(\mathrm{A}, \leq)$ is a poset.

Proof: Since $x \wedge x=x, x \leq x$ for all $x \in A$
Therefore $\leq$ is reflexive.
Suppose that $\mathrm{x}, \mathrm{y}, \mathrm{z} \in \mathrm{A}, \mathrm{x} \leq \mathrm{y}$ and $\mathrm{y} \leq \mathrm{z}$.
Then we have $y \wedge x=x \wedge y=x$ and $z \wedge y=y \wedge z=y$.
Now $x=x \wedge y=x \wedge y \wedge z=x \wedge z . \Rightarrow x \wedge z=z \wedge x=x$
Therefore, $\mathrm{x} \leq \mathrm{z}$. This shows that $\leq$ is transitive.
Suppose that $x, y \in A, x \leq y$ and $y \leq x \Rightarrow y \wedge x=x \wedge y=x$ and $y \wedge x=x \wedge y=y$.
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This shows that $\mathrm{x}=\mathrm{y}$. Therefore $\leq$ is anti-symmetric. Hence $(\mathrm{A}, \leq)$ is poset.
2. 1 Note: If $A$ is a Pre $A^{*}$-algebra with $1,0,2$ then $x \leq 1(x \wedge 1=1 \wedge x=x)$, for all $x \in A$ and $2 \leq x$ $(x \wedge 2=2 \wedge x=2)$. This shows that 1 is the greatest element and 2 is the least element of the poset. The Hasse diagram of the poset $(\mathrm{A}, \leq)$ is given by


## Diagram 2.1

We have $\mathrm{A} \times \mathrm{A}=\left\{\mathrm{a}_{1}=(1,1), \quad \mathrm{a}_{2}=(1,0), \quad \mathrm{a}_{3}=(1,2), \quad \mathrm{a}_{4}=(0,1), \quad \mathrm{a}_{5}=(0,0)\right.$, $\left.\mathrm{a}_{6}=(0,2), \mathrm{a}_{7}=(2,1), \mathrm{a}_{8}=(2,0), \mathrm{a}_{9}=(2,2)\right\}$ is a Pre $\mathrm{A}^{*}$-algebra under point wise operation and A $\times \mathrm{A}$ is having four central elements and remaining are non central elements, among that $\mathrm{a}_{9}=(2,2)$ is satisfying the property that $\mathrm{a}_{9}{ }^{\sim}=\mathrm{a}_{9}$. The Hasse diagram is of the poset $(\mathrm{A} \times \mathrm{A}, \leq)$ given below


Observe that, $x \leq a_{1}, x \wedge a_{1}=a_{1} \wedge x=x$ and $a_{9} \leq x\left(x \wedge a_{9}=a_{9} \wedge x=a_{9}\right)$ for all $x \in A \times A$. This shows that $a_{1}$ is the greatest element and $a_{9}$ is the least element of $A \times A$.

We have $\mathbf{2} \times \mathbf{3}=\left\{a_{1}=(1,1), a_{2}=(0,0), a_{3}=(1,0), a_{4}=(0,1), a_{5}=(2,2)\right.$,
$\left.a_{6}=(1,2)\right\}$ is a Pre $A^{*}$-algebra under point wise operation having four central elements, two noncentral elements and no element is satisfying the property that $\mathrm{a}^{\sim}=\mathrm{a}$.
The Hasse diagram for $(\mathbf{2} \times \mathbf{3}, \leq)$ as given below


Observe that, $x \leq a_{1}$, that is, $x \wedge a_{1}=a_{1} \wedge x=x$ and $a_{5} \leq x\left(x \wedge a_{5}=a_{5} \wedge x=a_{5}\right)$ for all $x \in \mathbf{2} \times \mathbf{3}$. This shows that $a_{1}$ is the greatest element and $a_{5}$ is the least element of $\mathbf{2} \times \mathbf{3}$.
2. 1. Theorem: In the partially ordered set $(A, \leq)$, for any $x \in A$, supremum of $\left\{x, x^{\sim}\right\}=x \vee x^{\sim}$ and infimum $\left\{x, x^{\sim}\right\}=x \wedge x^{\sim}$.

Proof: We have $\left(x \vee x^{\sim}\right) \wedge x=x$ and $x^{\sim} \wedge\left(x \vee x^{\sim}\right)=x^{\sim}$
Therefore $\mathrm{x} \leq \mathrm{x} \vee \mathrm{X}^{\sim}$ and $\mathrm{x}^{\sim} \leq \mathrm{x} \vee \mathrm{x}^{\sim}$
Hence $x \vee x^{\sim}$ is an upper bound of $\left\{x, x^{\sim}\right\}$
Suppose $n$ is an upper bound of $\left\{x, x^{\sim}\right\}$
That is, $\mathrm{x} \leq \mathrm{n}, \mathrm{x}^{\sim} \leq \mathrm{n} \Rightarrow \mathrm{n} \wedge \mathrm{x}=\mathrm{x}$, and $\mathrm{n} \wedge \mathrm{x}^{\sim}=\mathrm{x}^{\sim}$
Now $n \wedge\left(x \vee x^{\sim}\right)=(n \wedge x) \vee\left(n \wedge x^{\sim}\right)=x \vee x^{\sim}$
This shows that $\mathrm{x} \vee \mathrm{x}^{\sim} \leq \mathrm{n}$
Therefore $x \vee x^{\sim}$ is a least upper bound of $\left\{x, x^{\sim}\right\}$
This shows that supremum of $\left\{x, x^{\sim}\right\}=x \vee x^{\sim}$
Again we have $\left(x \wedge x^{\sim}\right) \wedge x=x \wedge x^{\sim}$ and $\left(x \wedge x^{\sim}\right) \wedge x^{\sim}=x \wedge x^{\sim}$
Therefore $\mathrm{x} \wedge \mathrm{x}^{\sim} \leq \mathrm{x}$ and $\mathrm{x} \wedge \mathrm{x}^{\sim} \leq \mathrm{x}^{\sim}$
Hence $x \wedge x^{\sim}$ is a lower bound of $\left\{x, x^{\sim}\right\}$
Suppose $m$ is a lower bound of $\left\{x, x^{\sim}\right\}$
That is, $m \leq x, m \leq x^{\sim} \Rightarrow m \wedge x=m$, and $m \wedge x^{\sim}=m$
Now $m \wedge\left(x \wedge x^{\sim}\right)=(m \wedge x) \wedge x^{\sim}=m \wedge x^{\sim}=m$
This shows that $m \leq x \wedge x^{\sim}$
Therefore $x \wedge x^{\sim}$ is a greatest lower bound of $\left\{x, x^{\sim}\right\}$
This shows that infimum of $\left\{x, x^{\sim}\right\}=x \wedge x^{\sim}$
2. 2. Theorem: In a poset $(A, \leq)$ with 1 , for any $x, y \in A, \operatorname{Inf}\{x, y\}=x \wedge y$.

Proof: We have $(x \wedge y) \wedge x=x \wedge y$ and $(x \wedge y) \wedge y=x \wedge y$
Therefore $x \wedge y \leq x$ and $x \wedge y \leq y$.
Hence $x \wedge y$ is a lower bound of $\{x, y\}$
Suppose $m$ is a lower bound of $\{x, y\}$
That is, $m \leq x, m \leq y \Rightarrow m \wedge x=m$ and $m \wedge y=m$
Now $m \wedge(x \wedge y)=(m \wedge x) \wedge y=m \wedge y=m$.
This shows that $m \leq x \wedge y$
Therefore $x \wedge y$ is a greatest lower bound of $\{x, y\}$
This shows that infimum of $\{x, y\}=x \wedge y$.
In general for a Pre $A^{*}$-algebra with 1 , $x \vee y$ need not be the l.u.b of $\{x, y\}$ in $(A, \leq)$. For example $2 \vee x=2 \wedge x=2, \forall x \in A$ is not a least upper bound. However we have the following theorem.
2. 3. Theorem: In a poset $(A, \leq)$ with 1 , for any $x, y \in B(A), \sup \{x, y\}=x \vee y$.

Proof: If $x, y \in B(A)$, then we have, $x \wedge(x \vee y)=x$ and $y \wedge(x \vee y)=y$
This shows that $x \leq x \vee y$ and $y \leq x \vee y$
Hence $x \vee y$ is an upper bound of $\{x, y\}$

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Suppose z is an upper bound of $\{\mathrm{x}, \mathrm{y}\}$, then $\mathrm{z} \wedge \mathrm{x}=\mathrm{x}, \mathrm{z} \wedge \mathrm{y}=\mathrm{y}$
Now $z \wedge(x \vee y)=(z \wedge x) \vee(z \wedge y)=x \vee y$
Therefore, $\mathrm{x} \vee \mathrm{y} \leq \mathrm{z}$.
Hence $\sup \{x, y\}=x \vee y$.
2.4 Theorem: In the poset $(A, \leq)$, if $x, y \in B(A)$, then $x \vee y \leq x \vee x^{\sim}$.

Proof: $\left(x \vee x^{\sim}\right) \wedge(x \vee y) \quad=\{x \wedge(x \vee y)\} \vee\left\{x^{\sim} \wedge(x \vee y)\right\}$

$$
=x \vee\left(x^{\sim} \wedge y\right)
$$

$$
=x \vee y
$$

Therefore $\mathrm{x} \vee \mathrm{y} \leq \mathrm{x} \vee \mathrm{x}^{\sim}$
2.5. Theorem: In the poset $(A, \leq)$, if $x \leq y$, then for any $z \in A$,
(a) $\mathrm{z} \wedge \mathrm{X} \leq \mathrm{z} \wedge \mathrm{y}$
(b) $\mathrm{z} \vee \mathrm{x} \leq \mathrm{z} \vee \mathrm{y}$

Proof: If $x \leq y$, then $x \wedge y=x$
(a) $(\mathrm{z} \wedge \mathrm{x}) \wedge(\mathrm{z} \wedge \mathrm{y})=\{(\mathrm{z} \wedge \mathrm{x}) \wedge \mathrm{z}\} \wedge \mathrm{y}=(\mathrm{z} \wedge \mathrm{x}) \wedge \mathrm{y}=\mathrm{z} \wedge \mathrm{x}$.

Therefore $\mathrm{z} \wedge \mathrm{x} \leq \mathrm{z} \wedge \mathrm{y}$
(b) $(\mathrm{z} \vee \mathrm{x}) \wedge(\mathrm{z} \vee \mathrm{y})=\mathrm{z} \vee(\mathrm{x} \wedge \mathrm{y})=\mathrm{z} \vee \mathrm{x}$

Therefore $\mathrm{z} \vee \mathrm{x} \leq \mathrm{z} \vee \mathrm{y}$
Now we are giving the following equivalent conditions for $x \leq y$.
2. 2. Lemma: In a Pre $A^{*}$-algebra (i) $x \leq y \Leftrightarrow x \wedge\left(x^{\sim} \vee y\right)=\left(x^{\sim} \vee y\right) \wedge x=x$
(ii) $x \leq y \Leftrightarrow y \wedge\left(y^{\sim} \vee x\right)=\left(y^{\sim} \vee x\right)=\left(y^{\sim} \vee x\right) \wedge y=x$

Proof: (i) If $x \leq y$

$$
\begin{aligned}
& \Leftrightarrow x \wedge y=x \\
& \Leftrightarrow x \wedge(x \sim y)=(x \sim \vee y) \wedge x=x
\end{aligned}
$$

(ii) If $x \leq y \Leftrightarrow y \wedge x=x$

$$
\Leftrightarrow y \wedge\left(y^{\sim} \vee x\right)=\left(y^{\sim} \vee x\right) \wedge y=x
$$

Now we prove modular type results in the following lemma.
2.3 Lemma: In the poset $(A, \leq)$, if $x \leq y \Rightarrow x \vee(y \wedge z)=y \wedge(x \vee z)$.

Proof: Suppose $x \leq y$ then $y \wedge x=x$
Now $y \wedge(x \vee z)=(y \wedge x) \vee(y \wedge z)=x \vee(y \wedge z)$
If $x, y \in B(A)$ then by theorem 2. 3 , $\sup \{x, y\}=x \vee y$. In general $x \vee y$ need not be an upper bound of $\{x, y\}$ in poset $(A, \leq)$. If $x \vee y$ is an upper bound of $\{x, y\}$ in poset $(A, \leq)$, then $A$ becomes Boolean algebra. Now we have the following theorem.
2.6. Theorem: If $A$ is a Pre $A^{*}$-algebra and $x \wedge(x \vee y)=x$ for all $x, y \in A$ then $(A, \leq)$ is a lattice.

Proof: By Theorem 2.2, we have every pair of elements have g.l.b and if $x \wedge(x \vee y)=x$ for all $x, y \in A$, then by theorem 2.3 we have every pair of elements have l.u.b. Hence $(\mathrm{A}, \leq)$ is a lattice.

Now we present an equivalent condition for a Pre $A^{*}$-algebra become a Boolean algebra.
2.7. Theorem: The following conditions are equivalent for any Pre $A^{*}$-algebra $\left(A, \wedge, \vee,(-)^{\sim}\right)$.
(1) A is a Boolean Algebra
(2) $x \leq x \vee y$ for all $x, y \in A$
(3) $y \leq x \vee y$ for all $x, y \in A$
(4) $x \vee y$ is an upper bound of $\{x, y\}$ in $(A, \leq)$ for all $x, y \in A$
(5) $x \vee y$ is an supremum of $\{x, y\}$ in $(A, \leq)$ for all $x, y \in A$
(6) $\mathrm{x} \vee \mathrm{x}^{\sim}$ is the greatest element in $(\mathrm{A}, \leq)$ for every $\mathrm{x} \in \mathrm{A}$

Proof: (1) $\Rightarrow$ (2) Suppose A be a Boolean algebra
Now $x \wedge(x \vee y)=x$ (by absorption law)
Hence $x \leq x \vee y$.
(2) $\Rightarrow$ (3) suppose $x \leq x \vee y$ then $x \wedge(x \vee y)=x$

Now $y \wedge(x \vee y)=y$. Therefore $y \leq x \vee y$.
(3) $\Rightarrow$ (4) Suppose that $y \leq x \vee y \Rightarrow y \wedge(x \vee y)=y$

Since $y \leq x \vee y$ then $x \vee y$ is upper bound of $y$
Now $x \wedge(x \vee y)=x$ (by supposition)
Therefore $x \leq x \vee y \Rightarrow x \vee y$ is upper bound of $x$
Hence $x \vee y$ is an upper bound of $\{x, y\}$.
(4) $\Rightarrow \mathbf{( 5 )}$ suppose $x \vee y$ is an upper bound of $\{x, y\}$

Suppose z is an upper bound of $\{x, y\}$, then $\mathrm{x} \leq \mathrm{z}, \mathrm{y} \leq \mathrm{z}$ that is $\mathrm{x} \wedge \mathrm{z}=\mathrm{x}, \quad \mathrm{y} \wedge \mathrm{z}=\mathrm{y}$
Now $z \wedge(x \vee y)=(z \wedge x) \vee(z \wedge y)=x \vee y$
Therefore $\mathrm{X} \vee \mathrm{y} \leq \mathrm{Z}$. Hence $\sup \{\mathrm{x}, \mathrm{y}\}=\mathrm{x} \vee \mathrm{y}$.
(5) $\Rightarrow$ (6) suppose $\sup \{\mathrm{x}, \mathrm{y}\}=\mathrm{x} \vee \mathrm{y}$ then $\mathrm{x}, \mathrm{y} \in B(A)$

Now $\sup \left\{x \vee x^{\sim}, y\right\}=x \vee x^{\sim} \vee y=x \vee x^{\sim}$
$\Rightarrow \mathrm{y} \leq \mathrm{x} \vee \mathrm{x}^{\sim}$
Therefore $\mathrm{x} \vee \mathrm{x}^{\sim}$ is the greatest element in $(\mathrm{A}, \leq)$.
(6) $\Rightarrow$ (1) suppose $x \vee x^{\sim}$ is the greatest element in $A$ then $y \leq x \vee x^{\sim}$
$\Rightarrow\left(\mathrm{x} \vee \mathrm{x}^{\sim}\right) \wedge \mathrm{y}=\mathrm{y}$
Now $y \vee(x \wedge y)=\left[\left(x \vee x^{\sim} \wedge y\right] \vee(x \wedge y)=\left[\left(x \vee x^{\sim}\right) \vee x\right] \wedge y\right.$
$=\left(x \vee x^{\sim}\right) \wedge y=y \quad$ (by supposition)
Therefore absorption law holds hence A is a Boolean algebra.
2.8. Theorem: Let $A$ be a pre $A^{*}$-algebra if $x \wedge x^{\sim}$ is the least element in
$(\mathrm{A}, \leq)$ for every $\mathrm{x} \in \mathrm{A}$, then A is a Boolean algebra.
Proof: Suppose $\mathrm{x} \wedge \mathrm{x}^{\sim}$ is the least element in $(\mathrm{A}, \leq)$ then $\mathrm{x} \wedge \mathrm{x}^{\sim} \leq \mathrm{y}$
$\Rightarrow\left(x \wedge x^{\sim}\right) \wedge y=x \wedge x^{\sim}$
Now $x \wedge(x \vee y)=\left[x \vee\left(x^{\sim} \wedge x\right)\right] \wedge(x \vee y)$
$=x \vee\left[\left(x \wedge x^{\sim}\right) \wedge y\right]$
$=\mathrm{x} \vee\left(\mathrm{x} \wedge \mathrm{x}^{\sim}\right)$ (by supposition)
$=\mathrm{x}$
Therefore $\mathrm{x} \wedge(\mathrm{x} \vee \mathrm{y})=\mathrm{x}$, absorption law holds.
Therefore A is a Boolean algebra.
2.9. Theorem: Let $A$ be a Pre $A^{*}$-algebra and $\mathrm{a} \in \mathrm{A}$. Let
$A_{a}=\{\mathrm{x} \in \mathrm{A} / \mathrm{a} \wedge \mathrm{x}=\mathrm{x}\}$.Then $A_{a}$ is closed under the operations $\wedge$ and $\vee$. Also for any $\mathrm{x} \in A_{a}$ define, $x^{a}=\mathrm{a} \wedge \mathrm{x}^{\sim}$. Then $\left(A_{a}, \wedge, \vee,^{\mathrm{a}}\right)$ is a Pre $\mathrm{A}^{*}$-algebra with 1 (here a is itself is the identity for $\wedge$ in $A_{a}$; that is 1 in $A_{a}$ ).

Proof: Let $x, y \in A_{a}$. Then $\mathrm{a} \wedge \mathrm{x}=\mathrm{x}$ and $\mathrm{a} \wedge \mathrm{y}=\mathrm{y}$.
Now $\mathrm{a} \wedge(\mathrm{x} \wedge \mathrm{y})=(\mathrm{a} \wedge \mathrm{x}) \wedge \mathrm{y}=\mathrm{x} \wedge \mathrm{y} \Rightarrow \mathrm{x} \wedge \mathrm{y} \in A_{a}$
Also $\mathrm{a} \wedge(\mathrm{x} \vee \mathrm{y})=(\mathrm{a} \wedge \mathrm{x}) \vee(\mathrm{a} \wedge \mathrm{y})=\mathrm{x} \vee \mathrm{y} \Rightarrow \mathrm{x} \vee \mathrm{y} \in A_{a}$
Therefore $A_{a}$ is closed under the operation $\wedge$ and $\vee$.
$\mathrm{a} \wedge x^{a}=\mathrm{a} \wedge\left(\mathrm{a} \wedge \mathrm{x}^{\sim}\right)=\mathrm{a} \wedge \mathrm{x}^{\sim}=x^{a} \Rightarrow x^{a} \in A_{a}$
Thus $A_{a}$ is closed under ${ }^{\text {a }}$.
Now for any $\mathrm{x}, \mathrm{y}, \mathrm{z} \in A_{a}$
(1) $x^{a a}=\left(a \wedge x^{\sim}\right)^{a}=a \wedge\left(a \wedge x^{\sim}\right)^{\sim}=a \wedge\left(a^{\sim} \vee x\right)=a \wedge x=x$
(2) $x \wedge x=(a \wedge x) \wedge(a \wedge x)=a \wedge x=x$
(3) $x \wedge y=(a \wedge x) \wedge(a \wedge y)=(a \wedge y) \wedge(a \wedge x)=y \wedge x$
(4) $(x \wedge y)^{a}=a \wedge(x \wedge y)^{\sim}=a \wedge\left(x \sim y^{\sim}\right)$

$$
\begin{aligned}
& =\left(a \wedge x^{\sim}\right) \vee\left(a \wedge y^{\sim}\right) \\
& =x^{a} \vee y^{b}
\end{aligned}
$$

(5) $x \wedge(y \wedge z)=(a \wedge x) \wedge\{(a \wedge y) \wedge(a \wedge z)\}$

$$
\begin{aligned}
& =a \wedge\{x \wedge(y \wedge z)\} \\
& =a \wedge\{(x \wedge y) \wedge z\}(\text { since } x, y, z \in A) \\
& =(x \wedge y) \wedge z
\end{aligned}
$$

(6) $x \wedge(y \vee z)=(a \wedge x) \wedge\{(a \wedge y) \vee(a \wedge z)\}$

$$
=\{(\mathrm{a} \wedge \mathrm{x}) \wedge(\mathrm{a} \wedge \mathrm{y})\} \vee\{(\mathrm{a} \wedge \mathrm{x}) \wedge(\mathrm{a} \wedge \mathrm{z})\}
$$

$$
=\{\mathrm{a} \wedge(\mathrm{x} \wedge \mathrm{y})\} \vee\{(\mathrm{a} \wedge(\mathrm{x} \wedge \mathrm{z})\}
$$

$$
=(x \wedge y) \vee(x \wedge z)
$$

(7) $x \wedge\left(x^{a} \vee y\right)=x \wedge\left\{\left(a \wedge x^{\sim}\right) \vee y\right\}$

$$
\begin{aligned}
& =\left\{\mathrm{x} \wedge\left(\mathrm{a} \wedge \mathrm{x}^{\sim}\right)\right\} \vee(\mathrm{x} \wedge \mathrm{y}) \\
& =\left(\mathrm{x} \wedge \mathrm{x}^{\sim}\right) \vee(\mathrm{x} \wedge \mathrm{y})(\text { since } \mathrm{a} \wedge \mathrm{x}=\mathrm{x})
\end{aligned}
$$

$$
\begin{aligned}
& =x \wedge(x \sim y) \\
& =x \wedge y
\end{aligned}
$$

Finally $\mathrm{x} \in A_{a}$ implies that $\mathrm{a} \wedge \mathrm{x}=\mathrm{x}=\mathrm{x} \wedge \mathrm{a}$. Thus $\left(A_{a}, \wedge, \vee{ }^{\mathrm{a}}{ }^{\prime}\right)$ is a Pre $\mathrm{A}^{*}$-algebra with a as identity for $\wedge$.
2.10. Theorem: Let $\mathrm{a}, \mathrm{b}$ be elements in a Pre $\mathrm{A}^{*}$-algebra A such that $a \leq b$. Then the following hold.
(1) $a \wedge b=a$
(2) The map $\alpha_{a, b}: A_{b} \rightarrow A_{a}$ defined by $\alpha_{a, b}(\mathrm{x})=\mathrm{a} \wedge \mathrm{x}$ for all $\mathrm{x} \in A_{b}$ is a homomorphism of Pre $\mathrm{A}^{*}$-algebras.
(3) $\alpha_{a, b}\left(\mathrm{~B}\left(A_{b}\right)\right) \subseteq \mathrm{B}\left(A_{a}\right)$
(4)If $\mathrm{a} \leq \mathrm{b} \leq \mathrm{c}$ then $\alpha_{a, b}$ o $\alpha_{b, c}=\alpha_{a, c}$
(5) $\alpha_{a, a}$ is the identity map on $A_{a}$

Proof: Suppose that $a \leq b$
(1) We have $a \leq b \Rightarrow a \wedge b=\mathrm{a}$
(2) Let $\mathrm{x}, \mathrm{y} \in A_{b}$.Then $\alpha_{a, b}(\mathrm{x} \wedge \mathrm{y})=\mathrm{a} \wedge(\mathrm{x} \wedge \mathrm{y})$

$$
\begin{aligned}
& =(\mathrm{a} \wedge \mathrm{x}) \wedge(\mathrm{a} \wedge \mathrm{y}) \\
& =\alpha_{a, b}(\mathrm{x}) \wedge \alpha_{a, b}
\end{aligned}
$$

and $\alpha_{a, b}(\mathrm{x} \vee \mathrm{y})=\mathrm{a} \wedge(\mathrm{x} \vee \mathrm{y})$

$$
\begin{aligned}
& =(\mathrm{a} \wedge \mathrm{x}) \vee(\mathrm{a} \wedge \mathrm{y}) \\
& =\alpha_{a, b}(\mathrm{x}) \vee \alpha_{a, b}
\end{aligned}
$$

Also $\alpha_{a, b}\left(\mathrm{x}^{\mathrm{b}}\right)=\mathrm{a} \wedge \mathrm{x}^{\mathrm{b}}$

$$
\begin{aligned}
& =a \wedge\left(b \wedge x^{\sim}\right) \\
= & (a \wedge b) \wedge x^{\sim} \\
= & a \wedge x^{\sim} \\
= & a \wedge\left(a^{\sim} \vee x^{2}\right) \\
= & a \wedge(a \wedge x)^{\sim} \\
= & (a \wedge x)^{a} \\
= & \left(\alpha_{a, b}(x)\right)^{a}
\end{aligned}
$$

Therefore $\alpha_{a, b}$ is a homomorphism of Pre $\mathrm{A}^{*}$-algebras.
(3) Let $\mathrm{x} \in \mathrm{B}\left(A_{b}\right)$.

Then $\mathrm{x} \vee \mathrm{x}^{\mathrm{b}}=\mathrm{b}$ (since b is identity in $\left.A_{b}\right)$ and therefore $\mathrm{b}=\mathrm{x} \vee\left(\mathrm{b} \wedge \mathrm{x}^{\sim}\right)$

$$
\begin{align*}
& \text { Now } \mathrm{b}=b \wedge b=\mathrm{b} \wedge\left(\mathrm{x} \vee\left(\mathrm{~b} \wedge \mathrm{x}^{\sim}\right)\right) \\
& =(b \wedge x) \vee\left(b \wedge x^{\sim}\right) \\
& =\mathrm{b} \wedge\left(\mathrm{x} \vee \mathrm{x}^{\sim}\right)  \tag{i}\\
& \text { Now } \alpha_{a, b}(\mathrm{x}) \vee\left[\alpha_{a, b}(\mathrm{x})\right]^{\mathrm{a}}=(\mathrm{a} \wedge \mathrm{x}) \vee(\mathrm{a} \wedge \mathrm{x})^{\mathrm{a}} \\
& =(a \wedge x) \vee\left[a \wedge(a \wedge x)^{\sim}\right] \\
& =(\mathrm{a} \wedge \mathrm{x}) \vee\left[\mathrm{a} \wedge\left(\mathrm{a}^{\sim} \vee \mathrm{x}^{\sim}\right)\right] \\
& =a \wedge\left[x \vee\left(a^{\sim} \vee x^{\sim}\right)\right] \\
& =a \wedge\left[a^{\sim} \vee\left(x \vee x^{\sim}\right)\right] \\
& =\mathrm{a} \wedge\left(x \vee x^{\sim}\right) \\
& =(\mathrm{a} \wedge \mathrm{~b}) \wedge\left(\mathrm{x} \vee \mathrm{x}^{\sim}\right) \\
& =\mathrm{a} \wedge\left[\mathrm{~b} \wedge\left(\mathrm{x} \vee \mathrm{x}^{\sim}\right)\right] \\
& =\mathrm{a} \wedge \mathrm{~b} \quad(\mathrm{by}(\mathrm{i})) \\
& =\mathrm{a} \text {, which is } 1 \text { in } A_{a}
\end{align*}
$$

Therefore $\alpha_{a, b}(\mathrm{x}) \in \mathrm{B}\left(A_{a}\right)$
Thus $\alpha_{a, b}\left(\mathrm{~B}\left(A_{b}\right)\right) \subseteq \mathrm{B}\left(A_{a}\right)$
(4)Let $\mathrm{a} \leq \mathrm{b} \leq \mathrm{c}$
$\left[\alpha_{a, b} \circ \alpha_{b, c}\right](\mathrm{x})=\alpha_{a, b}\left[\alpha_{b, c}(\mathrm{x})\right]$

$$
\begin{aligned}
& =\alpha_{a, b}[\mathrm{~b} \wedge \mathrm{x}] \\
& =\mathrm{a} \wedge \mathrm{~b} \wedge \mathrm{x} \\
& =\mathrm{a} \wedge \mathrm{x} \\
& =\alpha_{a, c}(\mathrm{x})
\end{aligned}
$$

Therefore $\alpha_{a, b}$ o $\alpha_{b, c}=\alpha_{a, c}$
(5) $\alpha_{a, a}(\mathrm{x})=\mathrm{a} \wedge \mathrm{x}=\mathrm{x}$ for all $\mathrm{x} \in A_{a}$

Then $\alpha_{a, a}$ is identity map on $A_{a}$.

## 3. Conclusion

This manuscript illustrates the nature of the Pre-A*-algebra like a partially ordered set. With respect to binary operation $\wedge$, defined a relation $\leq$ on a Pre- $A^{*}$-algebra and observed that such a Pre-A*-algebra as a partially ordered set with respect to the relation $\leq$ and derived corresponding results. It has been observed a necessary condition for a Pre- $\mathrm{A}^{*}$-algebra to become a lattice with respect to binary operation $\wedge$. For any $\mathrm{a} \in \mathrm{A}$ defined a set $A_{a}=\{\mathrm{x} \in \mathrm{A} / \mathrm{a} \wedge \mathrm{x}=\mathrm{x}\}$ and $x^{a}=\mathrm{a} \wedge \mathrm{x}^{\sim}$ , observed that $\left(A_{a}, \wedge, \vee,{ }^{\mathrm{a}}\right)$ is a Pre $\mathrm{A}^{*}$-algebra. Also by defining a mapping $\alpha_{a, b}$ from $A_{b}$ to $A_{a}$ by $\alpha_{a, b}(\mathrm{x})=\mathrm{a} \wedge \mathrm{x}$ for all $\mathrm{x} \in A_{b}$, confirmed a homomorphism of Pre $\mathrm{A}^{*}$-algebras.

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